

The Design of A Distributed Rating Scheme for Peer-to-peer Systems

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Outline

- Research motivations
- Basic design issues in P2P rating schemes
- A distributed rating scheme to incentivize cooperation in P2P file-sharing systems
- Dealing with collusion and malice
- Conclusion & future work

Research motivations

- Object: P2P file-sharing systems
 - ❑ Open social communities.
 - ❑ An explicit reputation layer was ignored in the original design.
- Goal: Build reputation in such systems
 - ❑ Incentive for user participation
 - free-riding phenomenon [Adar et al. 2000][Saroiu et al. 2001]
 - ❑ Isolation of malicious users
 - distribution of inauthentic files
 - propagation of virus or worms [VBS.Gnutella][Fizzer.Kazza]

P2P rating: basic design issues

- “Distributed” rating
 - following P2P design philosophy.
- “Efficient” rating
 - low cost to run and maintain this reputation system.
- “Collusion-proof” rating
 - Effectiveness.

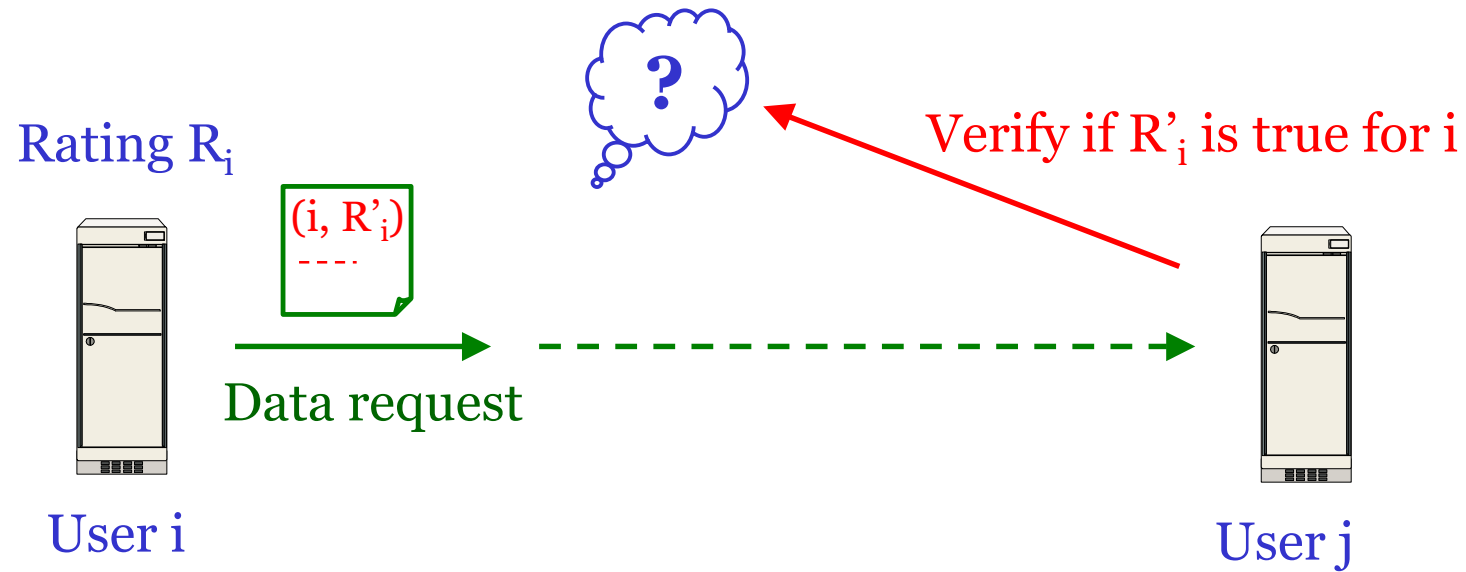
A distributed rating scheme

- To incentivize cooperation in P2P file-sharing systems
- Main components
 - ❑ Positive rating
 - ❑ Rating verification scheme

Positive rating

- The recognized service done to the community
 - R_i of user i : non-decreasing with the number of successful requests that it has satisfied within some sliding time window.
 - The higher R_i user i has, the better service it gets from the network.

Verification-based rating scheme



Two verification schemes

- Structured verification scheme (SVS)

- Each user has a set of designated supervisors which keep its up-to-date reputation information.

lightweight The supervisors are responsible for the verification.

- Unstructured verification scheme (UVS)

- A user j queries some of user i 's claimed customers for the verification, and believes i when the majority of the probed users reply with a "yes".

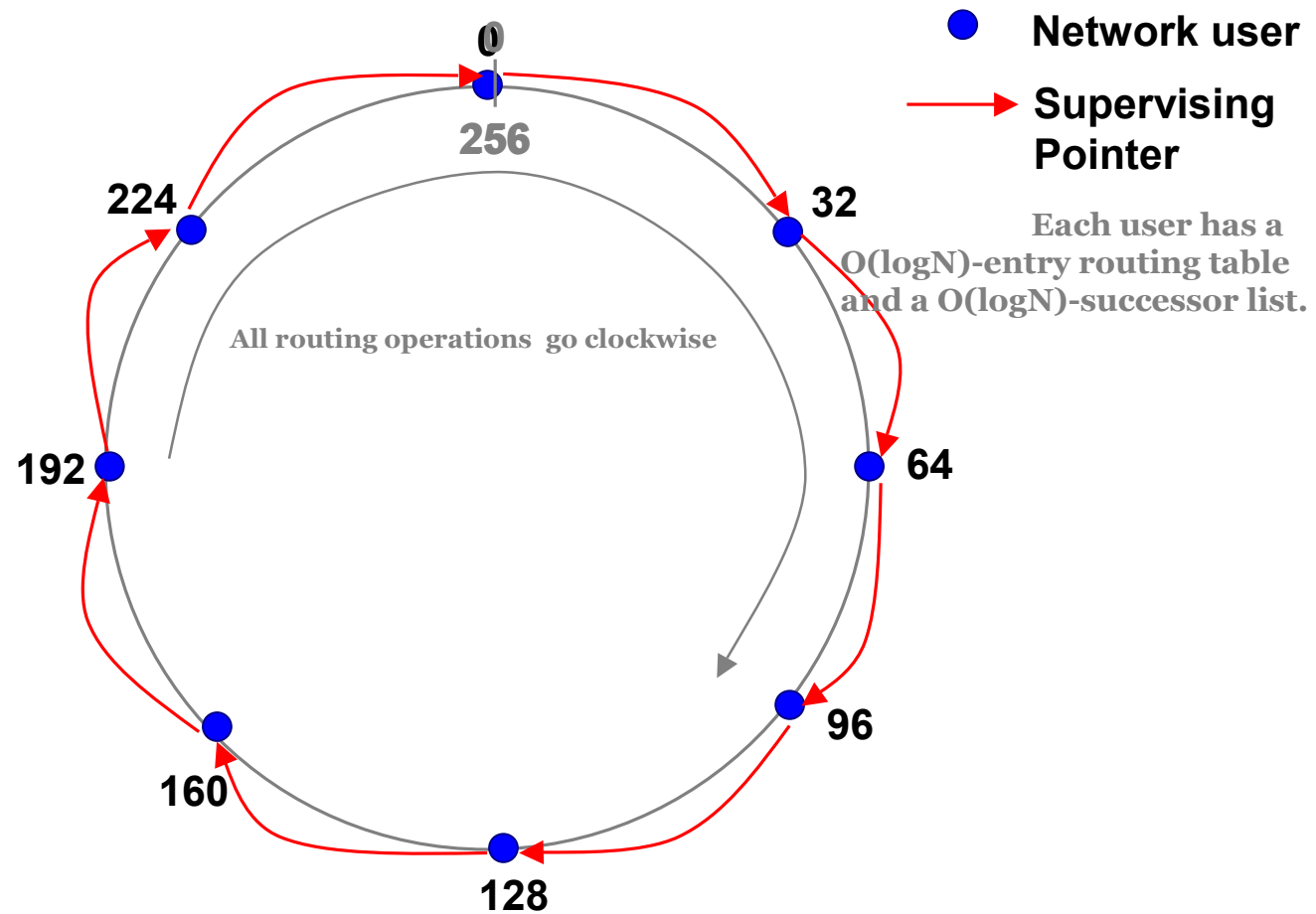
Assumption

- Users are distinguished by their IP addresses.
 - At a given time, one IP address corresponds to a unique user.

SVS – the supervising topology

- In the supervisory directed graph
 - ❑ Any user is random to its supervisors.
 - ❑ No small supervising loop exists.
 - ❑ There is a fast *reactive* approach for any user j to deliver a message to any other user i 's supervisors, and the path never includes i .

A Chord_[stoica2001] supervising overlay

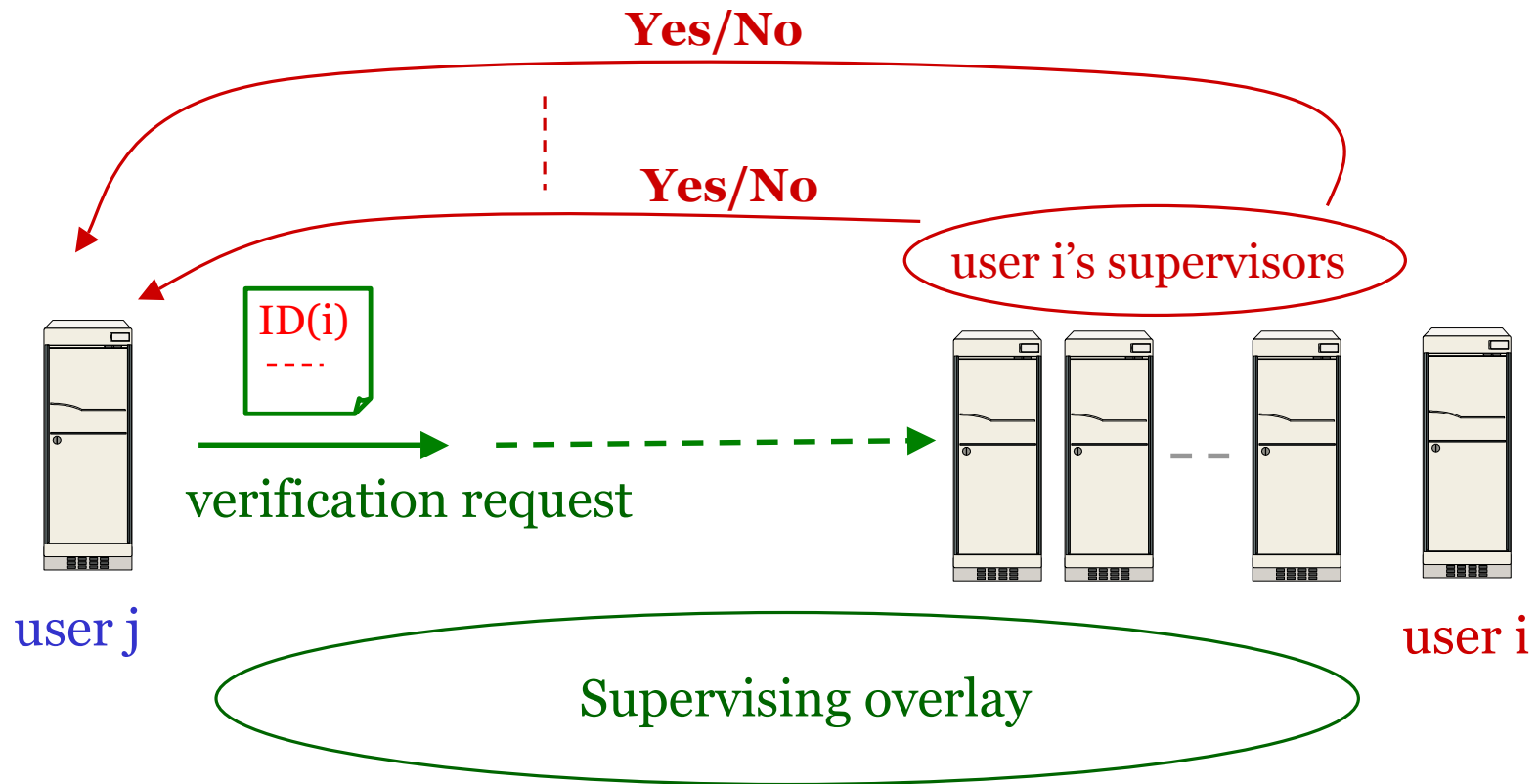


A Chord network with 8 users and 8-bit key space

SVS – the supervising topology

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Rating verification in SVS



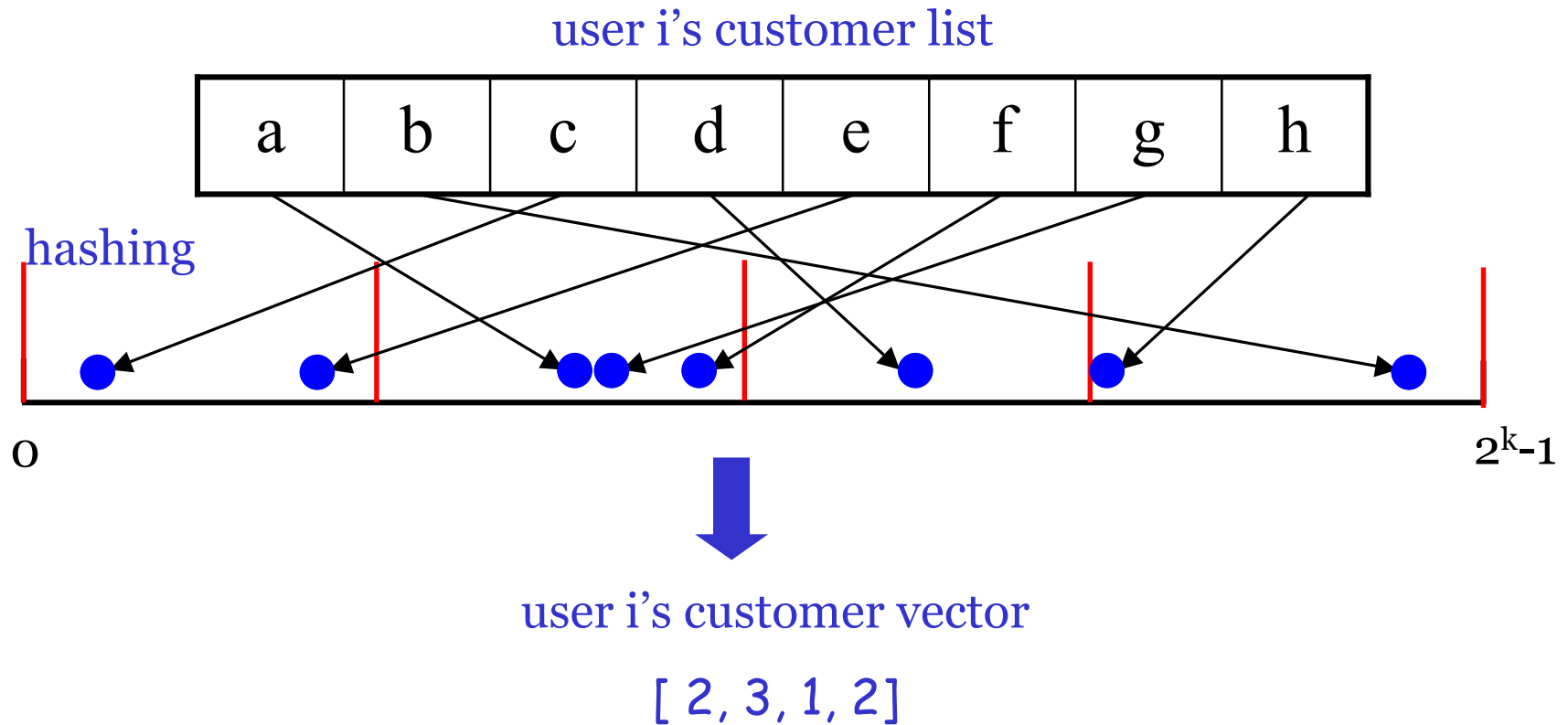
Structured verification scheme

- “Distributed” rating ✓
- “Collusion-proof” rating ✓
- “Efficient” rating ?
 - ❑ Extra cost to maintain a supervisory overlay when the underlying network is not DHT-based.
 - ❑ Repetitive actions when there are multiple supervisors.

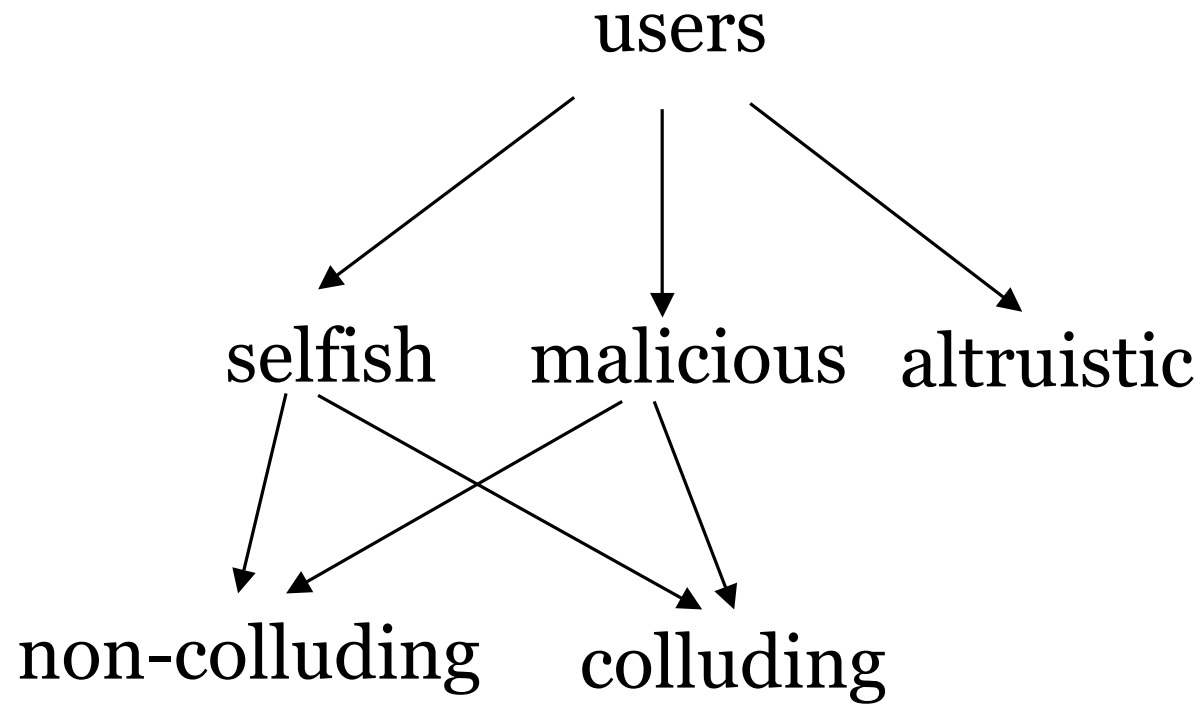
Unstructured verification scheme

1. When user j decides to verify user i 's rating, it gets a portion of i 's customer list, and asks the users on the list if i did the claimed service to them.
 - The customer samples should be random on the full customer list of node i .
 - Disclosure of full customer list could raise privacy concern, and incur high communication cost.
2. When the majority of the probed users reply with a “yes”, j is convinced that R'_i is R_i .

Randomly sampling without the complete customer list



P2P users



Colluding selfish users

- Possible solution 1: discrete rating
 - ❑ Grade★ : if a user has served no more than 10 users.
 - ❑ Grade★ ★ : if a user has served between 10 and 100 users.
 - ❑ Grade★ ★ ★ : if a user has served more than 1000 users.
- Possible solution 2: rating as virtual currency
 - ❑ A user has to pay (reduce its rating) for the service it claims to have received.
 - ❑ SVS: asks the requestor's supervisors for a payment.
 - ❑ LUVS: future work.

Colluding malicious users

- One possible strategy
 1. A user i quickly earns a high rating by faking transactions with other colluding users.
 2. User i then does bad things until earns bad-enough reputation.
 3. User i quits the network to clear its history,
 4. User i rejoins the network and repeats the above actions.

Conclusion & future work

- A simple distributed rating scheme to incentivize cooperation in P2P file-sharing systems.
 - Two distinct verification schemes
- Refine LUVS scheme to handle colluding selfish users.
- Refine our rating scheme to be collusion-proof.